Layered Area-Proportional Rectangle Contact Representations

Martin Nöllenburg, <u>Anaïs Villedieu</u>, Jules Wulms 16.9.2021 · GD 2021



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Background



Background















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CROWN problem (Contact Representation of Word Networks) Maximize $p = \sum_{e \in E_r} w_e$, where E_r are realised edge contacts



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Paths Cycles	[Barth et al. 2013]					



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Extension to many more graph classes
[Bekos et al. 2014]



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Extension to many more graph classes

[Bekos et al. 2014]

 \rightarrow based on generalized assignment problem approximation and graph decompositions

- Area proportional
 - Fixed aspect ratio
 - Unit height
 - \rightarrow fixed width



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- Contact representation
 - Rectangular dual
 - Border contacts



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Contact representations of planar graphs are a well-studied topic in graph theory, graph drawing and computational geometry. Vertices are represented by









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G = (V, E) layered graph with weight function $w : V \to \mathbb{R}^+$ Vertex weight \iff Rectangle width



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$R_{3,0}$	I	$R_{3,1}$.		$R_{3,2}$	I	$R_{3,3}$		$R_{3,4}$		$_5 R_{3,6}$	
	$R_{2,0}$		$R_{2,1}$	$R_{2,2}$		$R_{2,3}$	$R_{2,4}$		$R_{2,5}$		
	$R_{1,0}$			$R_{1,1}$		$R_{1,}$	2	$R_{1,3}$	$R_{1,4}$		
_	$R_{0,0}$			R_0	,1	R_0),2	$R_{0,3}$	$R_{0,2}$	4	$R_{0,5}$

Contact maximization

$R_{3,0}$		$R_{3,1}$ F		$R_{3,2}$	$_2$ $R_{3,3}$			$R_{3,4}$		$R_{3,5}$	$R_{3,6}$		
		$R_{2,0}$ I		$R_{2,1}$	$R_{2,2}$	$R_{2,2}$		$R_{2,3}$	-	$R_{2,4}$		$R_{2,5}$	
		$R_{\rm c}$	$1,\!0$	$R_{1,1}$		ŀ	$R_{1,2}$		$R_{1,3}$		$R_{1,2}$	4	
$R_{0,0}$		R	0,1	R	0,2	$R_{0,}$	3	$R_{0,}$	4	$R_{0,}$	5		



Contact maximization Maximize the amount of edges in G realised as box contacts \iff Maximize p when $\forall e \in E, w(e) = 1$


Problem definition

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Area minimization

	Ra	3,0		$R_{3,1}$		$R_{3,2}$	2	1	$R_{3,3}$	-	$R_{3,4}$	$R_{3,5}$	$R_{3,6}$
		$R_{2,0}$		$R_{2,1}$	$R_{2,2}$	R_2	2,3	R_2	2,4	R	2,5		
		$R_{\rm c}$	$1,\!0$	$R_{1,1}$		F	$R_{1,2}$		$R_{1,3}$	$R_{1,4}$			
	R_{0}	0,0			0,1		R_{i}	0,2	$R_{0,3}$	-	$R_{0,4}$	$R_{0,}$	5

Problem definition

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Area minimization

Minimize the total gap width in the representation



Overview Contact maximization Area minimization Greedy algorithm for L = 2Extension to L = 3? Flow network ILP formulation for $L \geq 3$

Overview





Overview







A representation with false contacts can be maximal



A representation with false contacts can be maximal \rightarrow not valid



They are forbidden to preserve edge semantics























Compute the contact maximal representation at every step



Layer 1 fan



















 \rightarrow Inactive layer ends with a fan





Compute the contact maximal representation at every step \rightarrow 2 contacts at most per step, 1 vertical, 1 horizontal



- Compute the contact maximal representation at every step
- \rightarrow 2 contacts at most per step, 1 vertical, 1 horizontal
- \rightarrow 3 rules when 2 contacts are not immediate

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- \rightarrow 2 contacts at most per step, 1 vertical, 1 horizontal
- \rightarrow 3 rules when 2 contacts are not immediate



Compute the contact maximal representation at every step

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Put the new rectangle as far left as possible





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Shift is strictly better new position is kept $\rightarrow k + 2$ contacts

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Shift is strictly better new position is kept $\rightarrow k + 2$ contacts

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6 - 26

Compute the contact maximal representation at every step \rightarrow 2 contacts at most per step, 1 vertical, 1 horizontal

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6 - 27

Compute the contact maximal representation at every step \rightarrow 2 contacts at most per step, 1 vertical, 1 horizontal

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Contact maximization

Compute the contact maximal representation at every step \rightarrow 2 contacts at most per step, 1 vertical, 1 horizontal

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Theorem 1 :

We can maximize the number of realized contacts in O(n) time for L = 2.













































Martin Nöllenburg, Anaïs Villedieu, Jules Wulms · Row-based Rectangle Contact Representations of Triangular Grid Graphs

Overview





Overview • |||| Contact maximization Area minimization Greedy algorithm for L = 2Extension to L = 3? Flow network ILP formulation for $L \geq 3$



Hard constraints

Hard constraints

Prevent overlap



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Hard constraints

- Prevent overlap
- Prevent false adjacencies





Hard constraints

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- Prevent false adjacencies



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Hard constraints

- Prevent overlap
- Prevent false adjacencies



Hard constraints

- Prevent overlap
- Prevent false adjacencies

Soft constraints

Increase counter with a failed horizontal contact



• ||||

Hard constraints

- Prevent overlap
- Prevent false adjacencies

Soft constraints

Increase counter with a failed horizontal contact

$$\begin{array}{c} x_{i,j} + w_{i,j} \\ \hline R_{i,j} \\ \hline R_{i+1,j} \\ \hline x_{i+1,j} \end{array}$$

Hard constraints

- Prevent overlap
- Prevent false adjacencies

- Increase counter with a failed horizontal contact
- Increase counter with a vertical contact



Hard constraints

- Prevent overlap
- Prevent false adjacencies

- Increase counter with a failed horizontal contact
- Increase counter with a vertical contact



Hard constraints

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Hard constraints

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Overview • |||| Area minimization Contact maximization Greedy algorithm for L = 2Extension to L = 3? Flow network ILP formulation for $L \geq 3$

Overview





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A representation with false adjacencies can have no gap width



Area minimization

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A representation with false adjacencies can have no gap width



A representation with false adjacencies can have no gap width but lose almost all contacts



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Area minimization

A representation with false adjacencies can have no gap width but lose almost all contacts



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Flow represents desired contacts between rectangles





Flow represents desired contacts between rectangles



Flow represents desired contacts between rectangles
 Buffer for excess flow



- Flow represents desired contacts between rectangles
- Buffer for excess flow
- Rectangles for flow to gaps





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Theorem 2 : We can minimize the total gap width in polynomial time.

Overview





Overview Contact maximization Area minimization Greedy algorithm for L = 2Extension to L = 3? Flow network ILP formulation for $L \geq 3$

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Layered wordles

Contact maximization on L=2



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Layered wordles

- Contact maximization on L=2
- Area minimization





Layered wordles

- \blacksquare Contact maximization on L=2
- Area minimization
- ILP model





Layered wordles

- Contact maximization on L=2
- Area minimization
- ILP model

Open problems



 \blacksquare Complexity for $L\geq 3$ problem



Layered wordles

- Contact maximization on L=2
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Open problems



Complexity for $L \ge 3$ problem
Algorithm for L = 3



Layered wordles

- Contact maximization on L=2
- Area minimization
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Open problems



- \blacksquare Complexity for $L\geq 3$ problem
- Algorithm for L = 3
- Generate layered graphs



Layered wordles

- $\blacksquare \quad {\rm Contact} \ {\rm maximization} \ {\rm on} \ L=2$
- Area minimization
- ILP model

Open problems



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- Algorithm for L = 3
- Generate layered graphs

vertex two representation all rectangles network clouds Representations length grid number edge ILP while f slidinggap realized algorithm range C optimal layers problem cost total between left see far triangular vertical far only adjacencies layer each words graphs DIOCK horizontal constraints forbidden vertices semantic solution